



# Applying Symbolic Execution to Test Implementations of a Network Protocol Against its Specification

**Hooman Asadian**, Paul Fiterau-Brostean, Bengt Jonsson,  
Konstantinos Sagonas





# Introduction

- Testing correctness of network protocol implementations is **essential**
- A successful software testing technique is **symbolic execution**
  - ⇒ However, it is not so effective at testing stateful systems

## This work:

- Presents a methodology that makes symbolic execution effective in
  - Testing network protocol implementations, and
  - Exposing requirement violations using **assumptions** and **assertions**
- Applies this methodology to implementations of the DTLS protocol
  - Revealing numerous new security vulnerabilities and bugs in them





# Methodology

1. Extract Specification Requirements
  - Represent the requirements by logical formulas
2. Augment the SUT with assumptions and assertions
  - Assume inputs under which a requirement can be violated
  - Assert that no forbidden action is performed
3. Symbolic Execution
  - Explores the paths in the augmented SUT
4. Test Case Construction and Validation
  - Confirm the bug on the unmodified SUT



# 1- Extract Specification Requirements

- Requirements from the protocol RFC are identified by particular keywords:
  - MUST, MUST NOT, SHOULD, SHOULD NOT, ...
- Two types of requirements are extracted:
  - Input validity requirements
  - Input-output requirements
- Represent the requirements by logical formulas





# Input Validity Requirements

- E.g., the DTLS 1.2 RFC states:

*“For each received record, **the receiver MUST verify** that the record contains **a sequence number that does not duplicate** the sequence number of any other record received during the life of this session.”*

- For a set of Records  $R$ , received during a DTLS session:

$$\forall r, r' \in R: r \neq r' \Rightarrow r.sequence\_number \neq r'.sequence\_number$$

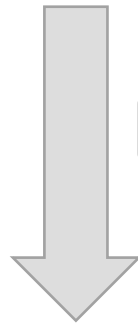




## 2- Augment the SUT with Assumptions

$\forall r, r' \in R: r \neq r' \Rightarrow r.sequence\_number \neq r'.sequence\_number$

$R = \{CH2, CKE, CCS\}$



Pair-wise conjunction

$assume (! ( CH2.sequence\_number \neq CKE.sequence\_number \& \\ CH2.sequence\_number \neq CCS.sequence\_number \& \\ CKE.sequence\_number \neq CCS.sequence\_number ))$





## 2- Augment the SUT with Assertions

- Add an assert statement to check if the implementation of the protocol uses invalid input in some forbidden way
- E.g., the DTLS 1.2 RFC:

*“Invalid records SHOULD be silently discarded ...”*

- Check whether progress occurs after reception of invalid records
  - Approximate this by successful completion of protocol interaction
  - Add failing assertion





## 3- Symbolic Execution

- Exploring the paths in the augmented SUT looking for assertion violation, crashes, memory errors, etc
- To achieve scalability:
  - Only make symbolic the relevant fields in a requirement
  - Other fields are given concrete values from a pre-captured session
  - Check one requirement at a time
- To ensure deterministic execution of the SUT:
  - De-randomize the SUT





# 4- Test Case Construction and Validation

- For each path, the tool returns:
  - A tuple of values for the symbolic fields
- For the sequence number experiment, we will have concrete values for *sequence\_number* in the participating records
- For concrete values that cause bugs:
  - Assign concrete values to relevant fields
  - Validate the bug by running the resulting test cases on the unmodified SUT





# Implementation and Application to DTLS

- Used KLEE as the symbolic execution engine
- Built a test harness that:
  - Captures the records a client and server exchange during a session
  - Is used to symbolically execute the SUT in order to check each requirement
- We implemented a shared library to facilitate test harness construction. It contains:
  - Helper functions
  - DTLS packet parser
  - Functions to make specific fields of records symbolic and to form Boolean expression in *assumes* and *asserts*

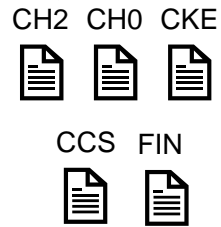




1 - Load the records from files to DTLS structured variable

4 - Feed the records to the side we are testing

### DTLS Test Harness



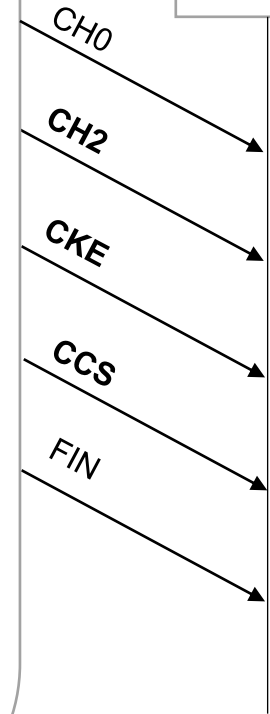
2- Make the relevant fields in the records symbolic

3 - Assume the negation of the requirement

### Shared library

```
make_symbolic (CH2.sequence_number)  
make_symbolic (CKE.sequence_number)  
make_symbolic (CCS.sequence_number)  
  
assume (!(CH2.sequence_number != CKE.sequence_number &  
CH2.sequence_number != CCS.sequence_number &  
CKE.sequence_number != CCS.sequence_number))
```

Server



assert(0)



# Evaluation

- We tested 4 DTLS libraries against 16 requirements:
  - 36 unique bugs
  - 7 vulnerabilities of which 6 are new

|                 | OpenSSL |               | Mbed TLS<br>2.22.0 | TinyDTLS <sup>E</sup> |         | TinyDTLS <sup>C</sup><br>53a0d97 |
|-----------------|---------|---------------|--------------------|-----------------------|---------|----------------------------------|
|                 | 1.0.1f  | 3.0.0-alpha12 |                    | 7068882               | 94205ff |                                  |
| Vulnerability   | 1       | 1             | –                  | 3                     | 3       | 2                                |
| Other           | –       | –             | –                  | 3                     | 4       | 1                                |
| Non-conformance | 2       | 2             | 3                  | 9                     | 10      | 10                               |





# TinyDTLS Reassembly Bug

- The DTLS 1.2 RFC specifies:

*“When a DTLS implementation receives a handshake message fragment, it MUST buffer it until it has the entire message”*

- Memory over-read when client/server reassemble a fragmented message
  - Occurs if the fragment length field is greater than the size of the actual fragment
- **Three** pull request attempts before the bug was fixed





# KLEE Experiences

- Protocol implementations define incoming/outgoing buffers sizes with respect to the Maximum Transmission Unit (MTU)
  - Memory over-read/over-write bugs can be missed by KLEE
  - **Our solution:** Allocate memory dynamically with respect to the size of the actual packets
  
- Significant interpretation slowdown when functions in cryptographic libraries are executed
  - Even in the absence of symbolic variables
  - Provided a benchmark in issue #1255 (**700% slowdown**)
  - **(Partial) solution:** Execute the functions as an external call



# Conclusion

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# Thank You for Listening

Replication materials available at:

<https://zenodo.org/record/5929867#.YkS3HSjMJJaT>

